Necessary condition for incompatibility of observables in general probabilistic theories

Sergey N. Filippov*

Institute of Physics and Technology, Russian Academy of Sciences, Moscow, Russia and Moscow Institute of Physics and Technology, Dolgoprudny, Moscow Region, Russia

Teiko Heinosaari[†] and Leevi Leppäjärvi[‡]

Turku Centre for Quantum Physics, Department of Physics and Astronomy, University of Turku, 20014 Turku, Finland (Received 7 October 2016; published 22 March 2017)

We quantify the intrinsic noise content of an observable in a general probabilistic theory and derive a noise content inequality for incompatible observables. We apply the derived inequality to standard quantum theory, the quantum theory of processes, and polytope state spaces. The noise content for positive operator-valued measures takes a particularly simple form and equals the sum of minimal eigenvalues of all the effects. We illustrate our findings with a number of examples including the introduced notion of reverse observables.

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I. INTRODUCTION

Quantum theory can be considered a particular instance within a wide range of probabilistic theories [1,2]. On the one hand, quantum theory inherits the general properties of probabilistic theories and, consequently, one may deduce some features already from a general operational framework. For instance, the limitations on broadcastable subsets of states can be derived from this generality [3]. On the other hand, particular properties of quantum theory, like specific constraints on nonlocality, partially fix its position with respect to other probabilistic theories [4]. As a result, specification of information-theoretic axioms may be sufficient for quantum theory to be derived [5].

A general probabilistic theory operates with notions of states and observables. The set of states S is convex since any probabilistic mixture of states must be a valid state. Observables are then affine functionals from the set of states ${\cal S}$ to the set of probability distributions. In standard quantum theory, states are associated with density operators, whereas observables are mathematically described by positive operator-valued measures (POVMs) [6–8]. However, when we are testing a quantum process, quantum channels are the examined objects and they are hence regarded as states, whereas observables can be described by process POVMs [9–11]. Theories describing the Popescu-Rohrlich (PR) box [12] and polytope state spaces serve as other examples of general probabilistic theories [13,14].

A set of observables in a general probabilistic theory may possess the property of being incompatible, which means that those observables cannot be seen as components of a single observable [15–18]. Incompatibility is a nonclassical feature, since in a general probabilistic theory with a classical state space all observables are compatible, while every nonclassical theory possesses some incompatible observables [19,20]. It is possible to compare the incompatibility of finite sets of observables in different probabilistic physical theories in a

reviewed. In Sec. III the noise content in observables is defined, and a sufficient condition for compatibility of a set of observables is formulated. The usage of the general condition

quantitative way [21–23]. Interestingly, quantum theory contains maximally incompatible pairs of observables, but only

when the underlying Hilbert space is infinite dimensional [24].

general probabilistic theories. The main goal of the present

investigation is to quantify the noise content for observables

in general probabilistic theories and to exploit it in deriving

a sufficient condition for compatibility, i.e., a necessary

condition for incompatibility for a collection of observables.

To demonstrate that the derived condition is noteworthy, we

use it to formulate a readily verifiable necessary condition for

incompatibility in quantum theory. To anticipate this result, the

condition takes the following form for POVMs: If m POVMs

are incompatible, then the sum of minimal eigenvalues of all

their elements is less than m-1. We illustrate our findings by

a number of examples including a class of reverse observables.

Consideration of standard quantum theory is followed by

theories with quantum processes as states, as well as the square

incompatibility has been investigated in several recent works

[25–28]. The conditions found in those works are tighter than

the condition presented in this work, but this is due to the fact

that they are applicable only for POVMs with some specific

structure or symmetry. Moreover, in contrast to most of the

earlier studies (see, for example, [15,17,22,29]), we do not

add noise to given observables but rather look for the intrinsic

We note that, in the case of POVMs, noise robustness of

This work focuses on the incompatibility of observables in

is then demonstrated in Sec. IV.

bit state space.

II. INCOMPATIBILITY OF OBSERVABLES IN GENERAL PROBABILISTIC THEORIES

A. States, effects, and observables

We begin by recalling the basic elements of the standard framework of general probabilistic theories (see, e.g., [30,31]

noise which is already present. We show that a meaningful nontrivial noise inequality can be derived already at the level of a general probabilistic theory. The paper is organized as follows. In Sec. II the incompatibility of observables in general probabilistic theories is

^{*}sergey.filippov@phystech.edu

[†]teiko.heinosaari@utu.fi

[‡]leille@utu.fi

for more detailed presentations). In a general probabilistic theory, the set of states S is a convex subset of a finite-dimensional real vector space V. The convexity is a result of the probabilistic nature of the theory, meaning that the convex sum $ps_1 + (1 - p)s_2$ is a state whenever s_1, s_2 are states and $0 \le p \le 1$.

We denote by F(S) the linear space of all affine functionals from S to \mathbb{R} ; i.e., a functional $e: S \to \mathbb{R}$ is in F(S) if it satisfies

$$e[ps_1 + (1-p)s_2] = pe(s_1) + (1-p)e(s_2)$$

for all $s_1, s_2 \in \mathcal{S}, 0 \leq p \leq 1$. For two functionals $e, f \in F(\mathcal{S})$, we denote $e \leq f$ if $e(s) \leq f(s)$ for all $s \in \mathcal{S}$. We further denote by $u \in F(\mathcal{S})$ the unit map satisfying u(s) = 1 for all $s \in \mathcal{S}$. The set of effects on \mathcal{S} is defined as

$$\mathcal{E}(\mathcal{S}) = \{ e \in F(\mathcal{S}) : 0 \leqslant e \leqslant u \},\$$

i.e., it is the convex subset of those affine functionals e for which $0 \le e(s) \le 1$ for all $s \in \mathcal{S}$. The set of effects arising as functionals on states is a particular example of an effect algebra [30]. In particular, $\mathcal{E}(\mathcal{S})$ has a partially defined sum e+f, which is simply the functional addition of e and f defined whenever $e+f \le u$.

An observable with a finite number of outcomes is a function $A: x \mapsto A_x$ from a finite outcome set $X \subset \mathbb{Z}$ to $\mathcal{E}(\mathcal{S})$. The number $A_x(s)$ is interpreted as the probability of getting the outcome x in a measurement of the observable A when the system is in the state s. As we must have $\sum_{x \in X} A_x(s) = 1$ for all $s \in \mathcal{S}$, we have the normalization condition $\sum_{x \in X} A_x = u$. We denote the set of observables with an outcome set X by \mathcal{O}_X , and by \mathcal{O} the set of all observables with a finite number of outcomes.

A special type of observable is the *trivial observable* T, which is, such that for each outcome x, $T_x(s) = T_x(s')$ for all $s, s' \in S$. We denote the set of trivial observables by T. Since the outcome probabilities for a trivial observable are the same for all states, it does not provide any information on an input state.

In what follows we recall the two most important instances of general probabilistic theories: standard quantum theory and the quantum theory of processes.

Example 1: Quantum theory. Let S_q be the convex set of density operators ϱ on a Hilbert space \mathcal{H} . Then the set of effects $\mathcal{E}(S_q)$, defined as affine mappings on S_q , can be represented as $e(\varrho) = \operatorname{tr}[\varrho E]$ for all states ϱ , where E is a self-adjoint operator satisfying the operator inequalities $0 \le E \le 1$. This correspondence is one to one, so effects can be identified with these effect operators. With this identification, an observable $A: x \mapsto A_x$ with a finite outcome set X is a POVM satisfying $\sum_{x \in X} A_x = 1$. A trivial observable T is of the form $T_x = p_x 1$, where p_x is a probability distribution on X.

Example 2: Quantum theory of processes. We denote by $\mathcal{L}(\mathcal{H})$ the bounded linear operators on a Hilbert space \mathcal{H} . Let \mathcal{S}_p be the set of completely positive and trace-preserving maps $\Phi: \mathcal{L}(\mathcal{H}_A) \to \mathcal{L}(\mathcal{H}_B)$, called quantum channels or processes. Then the set of effects $\mathcal{E}(\mathcal{S}_p)$ can be represented as the set of operators M on $\mathcal{H}_A \otimes \mathcal{H}_B$ satisfying $0 \leqslant M \leqslant \varrho \otimes \mathbb{I}$ for some density operator ϱ on \mathcal{H}_A . This representation is given as $e(\Phi) = \operatorname{tr}[\Omega_{\Phi} M]$, where Ω_{Φ} is the Choi operator of Φ ,

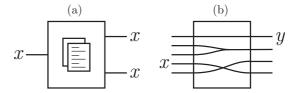


FIG. 1. (a) Action of the classical copying channel. (b) Example of the relabeling channel.

i.e., $\Omega_{\Phi} = (id \otimes \Phi)[|\psi_{+}\rangle\langle\psi_{+}|]$, where $\psi_{+} = \sum_{i=1}^{d} \phi_{i} \otimes \phi_{i}$ and $\{\phi_{i}\}_{i=1}^{d}$ is an orthonormal basis of \mathcal{H}_{A} . An important point is that this correspondence between affine maps and operators is not one to one; two operators M and M' correspond to the same effect e exactly when $M - M' = \omega \otimes \mathbb{1}$ for some traceless operator ω [32,33]. In this representation an observable $A: x \mapsto A_{x}$ with a finite outcome set X satisfies the normalization $\sum_{x \in X} A_{x} = \varrho \otimes \mathbb{1}$ for some density operator ϱ on \mathcal{H}_{A} . This kind of map is called a process POVM, or PPOVM for short [10]. A trivial PPOVM is of the form $T_{x} = p_{x}\xi_{x} \otimes \mathbb{1}$, where each ξ_{x} is a density operator on \mathcal{H}_{A} and p_{x} is a probability distribution. Two trivial PPOVMs $T_{x} = p_{x}\xi_{x} \otimes \mathbb{1}$ and $T'_{x} = p'_{x}\xi'_{x} \otimes \mathbb{1}$ correspond to the same trivial observable exactly when the probability distributions p_{x} and p'_{x} are the same.

B. Postprocessing of observables

A classical channel ν between outcome spaces X and Y is a right stochastic matrix with elements ν_{xy} , $x \in X$, $y \in Y$, i.e., $0 \le \nu_{xy} \le 1$ and $\sum_{y \in Y} \nu_{xy} = 1$. The number ν_{xy} is the transition probability for an element x to be transformed into y. Classical channels are often used to describe noise, but we can also think of a classical channel as an active transformation that is implemented on outcomes. In the following we recall two classes of classical channels that are used later.

Example 3: Copying the measurement outcomes. Measurement outcomes are just classical symbols and thus can be copied. To see copying as a classical channel, let $Y = X \times X$. The stochastic matrix v_{xy}^c related to copying is defined as $v_{xy}^c = 1$ if y = (x,x) and $v_{xy}^c = 0$ otherwise. This transforms any x to (x,x). Figure 1(a) depicts the action of the copying channel. Multiple applications of a copying channel allow one to make an arbitrary number of copies of an outcome x. If the number of copies equals m, then we call it an m-copying channel.

Example 4: Relabeling the measurement outcomes. The copying channels belong to a wider class of classical channels where measurement outcomes are relabeled deterministically into some other outcome. Let $f: X \to Y$ be a relabeling function. The derived stochastic matrix v_{xy}^f is defined as $v_{xy}^f = 1$ if f(x) = y and $v_{xy}^f = 0$ otherwise. In contrast to the copying procedure, generally several outcomes can be relabeled into a single new outcome; see Fig. 1(b).

Let A be an observable with an outcome set X and let ν be a classical channel between X and some other outcome space Y. We denote by $\nu \circ A$ the new observable defined as

$$(\nu \circ \mathsf{A})_{y} = \sum_{x \in X} \nu_{xy} \mathsf{A}_{x} \tag{1}$$

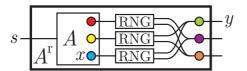


FIG. 2. A^r is the reverse observable with respect to A. Outcome y of observable A does not contribute to the outcome y of observable A^r, so they are illustrated by complementary colors. RNG stands for a random number generator which uniformly chooses outcome $y \neq x$.

for all outcomes $y \in Y$. Physically, the observable $v \circ A$ is implemented by first measuring A and then using the classical channel v on each obtained measurement outcome. This way of forming new observables gives rise to a preorder in the set of observables [34–36]. Namely, for two observables A and B, we say that B is a *postprocessing* of A if there exists a classical channel v such that $B = v \circ A$.

Example 5: Reverse observable. A reversing channel is a classical channel $\nu^r: X \to X$ such that $\nu^r_{xy} = 0$ if x = y and $\nu^r_{xy} = \nu^r_{x'y}$ for all $x, x' \neq y$. If the outcome set X contains N elements, then $\nu^r_{xy} = \nu^r_{x'y} = \frac{1}{N-1}$ for all $x, x' \neq y$. For each observable A, the observable $A^r = \nu^r \circ A$ is called the reverse version of A. If A has N outcomes, then the reverse observable A^r takes the form

$$A_x^{\rm r} = \frac{1}{N-1} \sum_{y \neq x} A_y = \frac{1}{N-1} (u - A_x).$$
 (2)

The physical meaning of A^r is illustrated in Fig. 2. After A has been measured and outcome x has been obtained, we roll a fair dice with N-1 sides and randomly choose any outcome y different from x. This is taken to be the outcome of the new observable A^r , which is hence given by formula (2).

Example 6: Doubly reverse observable. Performing the reversing postprocessing two times, we get

$$A_y^{\rm rr} = \frac{1}{(N-1)^2} [A_y + (N-2)u], \tag{3}$$

or, concisely, $A^{rr} = (1 - \lambda)A + \lambda T$, where $\lambda = \frac{N(N-2)}{(N-1)^2}$ and T is the trivial observable with uniform distribution of outcomes. In the case of two outcomes (N=2), the doubly reverse observable coincides with the original one, i.e., $A^{rr} = A$.

As one would expect, a trivial observable T is a postprocessing of any other observable A. To see this, we define a classical channel v^T as $v_{xy}^T = T_y(s_0)$ for all x, where s_0 is any state. Then

$$(\nu^{\mathsf{T}} \circ \mathsf{A})_{y}(s) = \sum_{x \in X} \nu_{xy}^{\mathsf{T}} \mathsf{A}_{x}(s) = \sum_{x \in X} \mathsf{T}_{y}(s_{0}) \mathsf{A}_{x}(s) = \mathsf{T}_{y}(s_{0})$$
$$= \mathsf{T}_{y}(s),$$

showing that $\nu^T \circ A = T$. The classical channel ν^T just erases the outcome obtained in the A measurement and replaces it with a new outcome according to the measurement outcome distribution of T, which is the same for all states.

C. Incompatibility of observables

A collection of observables \mathcal{P} is *compatible* if there exists an observable C, with an outcome set Y, such that each

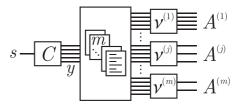


FIG. 3. Observables $A^{(1)}, \dots, A^{(m)}$ are compatible if each of them is a postprocessing of some observable C.

observable $A \in \mathcal{P}$ is a postprocessing of C. A compatible collection of observables can thus be implemented simultaneously by first measuring C, then copying the classical outcomes, and finally applying the relevant postprocessings to the copied outcomes. This definition is depicted in Fig. 3. If a set of observables is not compatible, then it is called *incompatible*.

Let $\{A^{(1)}, \ldots, A^{(m)}\}$ be a compatible set of m observables, with outcome sets $X^{(1)}, \ldots, X^{(m)}$, respectively. Thus, there exists an observable C and classical channels $\nu^{(1)}, \ldots, \nu^{(m)}$ such that

$$A^{(j)} = v^{(j)} \circ C, \quad j = 1, \dots, m.$$
 (4)

To see this definition of compatibility in an equivalent form, we denote

$$\mathsf{G}_{x^{(1)},\dots,x^{(m)}} = \sum_{y} \prod_{j=1}^{m} \nu_{yx^{(j)}}^{(j)} \mathsf{C}_{y} \tag{5}$$

for all $x^{(j)} \in X^{(j)}$, j = 1, ..., m. Then **G** is an observable, and from (4) it follows that

$$\mathsf{A}_{x^{(j)}}^{(j)} = \sum_{x^{(i)}: i \neq i} \mathsf{G}_{x^{(1)}, \dots, x^{(m)}}. \tag{6}$$

Thus, the compatibility of observables $A^{(1)}, \ldots, A^{(m)}$ implies that there exists a *joint observable* G with the outcome space $X^{(1)} \times \cdots \times X^{(m)}$ such that the observables are marginals of the joint observable. Conversely, starting from G and taking classical channels corresponding to relabeling functions that are projections, $pr_{\ell}: X^n \to X$, $pr_{\ell}(x_1, \ldots, x_n) = x_{\ell}$, we see that (6) is a special case of (4). As noted in [37] in the case of quantum observables, we conclude that a subset of observables is compatible if and only if they have a joint observable. The latter condition is usually taken as the definition of joint measurability of quantum observables [38].

III. NECESSARY CONDITION FOR INCOMPATIBILITY

A. Noise content of an observable

In order to formulate a necessary condition for incompatibility of observables, we first quantify their intrinsic fuzziness, or noise content, and then use the extraction of that noise in an explicit construction of a class of joint observables.

In a general probabilistic theory, one can introduce a procedure of mixing observables. Suppose $A: X \to \mathcal{E}(\mathcal{S})$ and $B: Y \to \mathcal{E}(\mathcal{S})$ are observables with outcome sets X and Y, respectively. Then a mixture of A and B, with a mixing parameter $0 \le t \le 1$, is an observable $C: X \cup Y \to \mathcal{E}(\mathcal{S})$ such that

$$C_z = tA_z + (1 - t)B_z \tag{7}$$

for all $z \in X \cup Y$, where A and B can be extended to $X \cup Y$ by defining $A_z = 0$ if $z \notin X$ and $B_z = 0$ if $z \notin Y$.

We are interested in a situation where one of the observables on the right-hand side of mixture (7) is not arbitrary but belongs to a some specified subset $\mathcal{N} \subseteq \mathcal{O}$ which describes noise in the measurement. If the target observable C is not in \mathcal{N} , then this requirement imposes limitations on possible values of the mixing parameter t.

For the following consideration, we fix a nonempty subset $\mathcal{N} \subseteq \mathcal{O}$ which describes noisy observables. Then, the physical meaning of Eq. (7) is to decompose an observable into its noisy part and the rest. A quantitative description of the noise content is attained by maximizing t. Therefore, for each observable A, we denote

$$w(A; \mathcal{N}) = \sup\{0 \le t \le 1 : tN + (1 - t)B = A$$
 for some $N \in \mathcal{N}$ and $B \in \mathcal{O}\}$ (8)

and call this quantity the *noise content of* A *with respect to* \mathcal{N} . We note that the observables N and B in (8) can be assumed to have the same outcome set as A.

Whenever $A_x \ge t N_x$ for some $0 \le t \le 1$ for all $x \in X$, we use the notation $A \geqslant N$. Suppose $0 \leqslant t < 1$ and $A \geqslant N$; then we can write A as a mixture

$$A = tN + (1 - t)\widetilde{A}, \tag{9}$$

where \widetilde{A} is the observable defined as

$$\widetilde{\mathbf{A}} = (1 - t)^{-1} (\mathbf{A} - t \mathbf{N}). \tag{10}$$

Conversely, if there exists some observable \widetilde{A} such that (9) holds, then $A \geqslant N$. Thus, one can reformulate the definition of noise content of A with respect to $\mathcal N$ as follows:

$$w(A; \mathcal{N}) = \sup\{0 \leqslant t \leqslant 1 : A \geqslant N \text{ for some } N \in \mathcal{N}\}.$$
 (11)

Specific properties of the map $A \mapsto w(A; \mathcal{N})$ depend on the choice of the subset \mathcal{N} . There are, however, some general features valid for any noise set \mathcal{N} . In particular, we observe the following:

- (a) If ν is a classical channel and $\nu \circ \mathcal{N} \subseteq \mathcal{N}$, then $w(\nu \circ \mathcal{N})$ $A; \mathcal{N} \geqslant w(A; \mathcal{N}).$
- (b) If \mathcal{N} is convex, then $w[sA + (1-s)B; \mathcal{N}] \geqslant$ $sw(A; \mathcal{N}) + (1 - s)w(B; \mathcal{N})$ for all $0 \le s \le 1$.

The first property follows directly from the definition of $w(A; \mathcal{N})$, while the latter is seen to be valid by first noticing

$$sA_x + (1-s)B_x \geqslant sw(A; \mathcal{N})N_x + (1-s)w(B; \mathcal{N})M_x$$

for some observables $N, M \in \mathcal{N}$ and all outcomes x. We denote

$$p_{\mathsf{A}} = sw(\mathsf{A}; \mathcal{N})/[sw(\mathsf{A}; \mathcal{N}) + (1-s)w(\mathsf{B}; \mathcal{N})],$$

$$p_{\mathsf{B}} = (1 - s)w(\mathsf{B}; \mathcal{N})/[sw(\mathsf{A}; \mathcal{N}) + (1 - s)w(\mathsf{B}; \mathcal{N})],$$

and then obtain

$$sA_x + (1-s)B_x$$

$$\geq [sw(A; \mathcal{N}) + (1-s)w(B; \mathcal{N})](p_AN + p_BM)_r$$

where $p_A N + p_B M \in \mathcal{N}$ as the set \mathcal{N} is convex.

The prototypical choice for \mathcal{N} is to take $\mathcal{N} = \mathcal{T}$, the set of all trivial observables. In this case, we simply say that w(A; T)

is the noise content of A. The set \mathcal{T} is convex and $\nu \circ \mathcal{T} \subseteq \mathcal{T}$ for all classical channels.

Proposition 1. Let A be an observable on a finite outcome

set X. Then $w(A; \mathcal{T}) = \sum_{x \in X} \inf_{s \in \mathcal{S}} A_x(s)$. Proof. Denote $a_x = \inf_{s \in \mathcal{S}} A_x(s)$ and $a = \sum_x a_x$. First assume that $a_x = 0$ for all $x \in X$ so that a = 0. Let $T \in T$ be a trivial observable and take any t, $0 \le t \le 1$, such that $A \geqslant T$. By our definitions this is equivalent to $A_x(s) \geqslant tT_x(s)$

for all
$$x \in X$$
 and $s \in S$, so that for all $x \in X$ we have that
$$0 = a_x = \inf_{s \in S} \mathsf{A}_x(s) \geqslant t \inf_{s \in S} \mathsf{T}_x(s) = t p_x,$$

where $p_x \equiv \mathsf{T}_x(s)$ is the probability distribution defined by T . Summing over x we get

$$0 = a = \sum_{x} a_x \geqslant t \sum_{x} p_x = t$$

 $0 = a = \sum_x a_x \geqslant t \sum_x p_x = t.$ Since also $0 \leqslant t \leqslant 1$, we must have t = 0, and since this holds for all $T \in \mathcal{T}$, by (11) we get that $w(A; \mathcal{T}) = a = 0$.

Second, assume that $a_x \neq 0$ at least for some $x \in X$. By similar arguments as above, we see that for all $x \in X$ we have $a_x \geqslant t' p_x'$, where $p_x' = \mathsf{T}_x'(s)$ is a probability distribution defined by some trivial observable $T' \in \mathcal{T}$ for some $0 \le t' \le 1$. Summing over all x we then get an upper bound for t' as $a = \sum_{x} a_{x} \ge t'$. We see that the upper bound is attained if we define T' as T'(s) = $p'_x = a_x/a$. Thus by (11) we have that $w(A; \mathcal{T}) = a.$

B. Joint measurement scheme

The joint measurement scheme that we next discuss is an elaboration of the one presented in [28]. The idea is that we first write the definition of compatibility in a slightly different way, then limit the defining conditions, and in this way we obtain a computable sufficient condition for compatibility.

From the definition, two observables A and B are compatible if there exists a third observable C and classical channels v_1 and v_2 such that $A = v_1 \circ C$ and $B = v_2 \circ C$. Let us consider a seemingly more general scheme, where we are asking for the existence of two observables C and D, classical channels v_1 , ν_2 , μ_1 , and μ_2 , and a mixing parameter t such that

$$A = t\nu_1 \circ C + (1 - t)\mu_1 \circ D, \tag{12}$$

$$B = t\nu_2 \circ C + (1 - t)\mu_2 \circ D.$$
 (13)

Thus, A and B are now required to be mixtures of postprocessings of C and D; see Fig. 4.

Clearly, conditions (12) and (13) reduce to the usual compatibility conditions when t = 1. Therefore, every compatible pair can be written in this new form. Conversely, if two observables A and B can be written in the form of conditions (12) and (13), then they are compatible. In fact, A and B are postprocessings of the mixed observable tC + (1 - t)D, but now the mixture has an extra outcome to keep track of which observable was measured each time. After measuring either C or D, we duplicate the outcome and postprocess with either v_1 and ν_2 or μ_1 and μ_2 , depending on the measured observable.

C. Incompatibility inequality

As a special case of the joint measurement scheme described previously, we limit the choice of classical channels

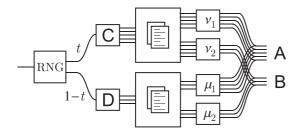


FIG. 4. The considered joint measurement scheme for two observables A and B consists of a random choice between two observables C and D, followed by separated postprocessings for both A and B that aim to approximate these observables.

 μ_1 and ν_2 to those that make observables $\mu_1 \circ D$ and $\nu_2 \circ C$ trivial. Since any trivial observable is a postprocessing of any other observable, we get all trivial observables, irrespective of C and D. Hence, conditions (12) and (13) reduce to

$$A = t \nu_1 \circ C + (1 - t) T_1,$$
 (14)

$$B = tT_2 + (1 - t)\mu_2 \circ D, \tag{15}$$

where T_1 and T_2 are arbitrary trivial observables. Since we have added an extra limitation to conditions (12) and (13), we cannot be sure anymore that a pair of compatible observables has this kind of representation. However, if $w(A; \mathcal{T}) \ge 1 - t$ and $w(B; \mathcal{T}) \ge t$, then by the definition of noise content we can find suitable observables C and D such that conditions (14) and (15) hold.

As a conclusion, we obtain the following result and its equivalent formulation.

Proposition 2. If A and B are two observables such that $w(A; T) + w(B; T) \ge 1$, then they are compatible.

Proposition 3. If A and B are incompatible observables, then w(A; T) + w(B; T) < 1.

The joint measurement scheme has a direct generalization for any finite number of observables. Let us consider m observables $A^{(1)}, \ldots, A^{(m-1)}$ and $A^{(m)}$. We can then generalize conditions (14) and (15) to

$$A^{(j)} = p_j v_j \circ C^{(j)} + (1 - p_j) T^{(j)}, \tag{16}$$

where $\mathsf{T}^{(j)}$ is an arbitrary trivial observable for each $j=1,\ldots,m$ and p_j is an arbitrary probability distribution. As above, if $w(\mathsf{A}^{(j)};\mathcal{T})\geqslant 1-p_j$ for all j we can make (16) hold. By summing over j we conclude the following generalization of Proposition 3.

Proposition 4. If $A^{(1)}, A^{(2)}, \dots, A^{(m)}$ are incompatible observables, then $w(A^{(1)}; \mathcal{T}) + \dots + w(A^{(m)}; \mathcal{T}) < m - 1$.

IV. APPLICATIONS OF THE INCOMPATIBILITY CONDITION

A. Eigenvalue condition for POVMs

If A is an observable in finite-dimensional quantum theory described by a POVM, we have that

$$\inf_{s \in \mathcal{S}} \mathsf{A}_{x}(s) = \min_{\psi \neq 0} \frac{\langle \psi \mid \mathsf{A}_{x} \psi \rangle}{\langle \psi \mid \psi \rangle}. \tag{17}$$

It follows that $\inf_{s \in \mathcal{S}} \mathsf{A}_x(s)$ is the smallest eigenvalue of the effect operator A_x . Hence, by Proposition 1 we conclude that $w(\mathsf{A}; \mathcal{T})$ is the sum of the minimal eigenvalues of operators A_x . Combining this with Proposition 4, we reach the following necessary condition for incompatibility.

Corollary 1. If $A^{(1)}, \ldots, A^{(m)}$ is a collection of m incompatible POVMs, then the sum of the minimal eigenvalues of all their effects is smaller than m-1.

We next illustrate the use of Corollary 1 in the case of reverse observables. Consider a *regular rank-1 POVM* A, i.e., the effects of A read $A_x = \frac{d}{N}P_x$, where d is the dimension of the Hilbert space, N is the number of outcomes, and P_x is a one-dimensional projection. Examples of regular rank-1 POVMs include all nondegenerate sharp POVMs and symmetric informationally complete POVMs.

As before, we denote by $A^r = v^r \circ A$ the reverse version of A. If A is a regular rank-1 POVM, then the smallest eigenvalue of each operator A_x^r is $\frac{N-d}{N(N-1)}$. Applying Corollary 1, we conclude that the reverse versions of m regular rank-1 POVMs with N outcomes are compatible if

$$N \geqslant (d-1)m+1. \tag{18}$$

It follows from this observation that, for instance, the reverse versions of two regular rank-1 POVMs in d=2 are compatible for all $N\geqslant 3$. One can readily find POVMs with two outcomes whose reverse versions are incompatible; this is the case whenever the original ones are incompatible since, in the case of two outcomes, reversing is a reversible classical channel. Since the reversing channel is more and more noisy when the number of outcomes increases, one may wonder if there are any incompatible collections of reverse POVMs when the number of outcomes is more than two. In the following example we present a triplet of regular rank-1 POVMs whose reverse versions are incompatible; the simple compatibility condition (18) is hence not trivial.

Example 7: Incompatible reverse POVMs. Consider three orthonormal bases $\{\varphi_i\}_{i=1}^3$, $\{\psi_i\}_{i=1}^3$, and $\{\chi_i\}_{i=1}^3$ in a three-dimensional Hilbert space \mathcal{H}_3 such that a set $\{\varphi_i, \psi_j, \chi_k\}$ is linearly independent for all fixed i, j, k. Let A, B, and C be the POVMs related to these bases, i.e., $A_i = |\varphi_i\rangle\langle\varphi_i|$, $B_i = |\psi_i\rangle\langle\psi_i|$, and $C_i = |\chi_i\rangle\langle\chi_i|$. The fact that the reverse POVMs A^r , B^r , and C^r are incompatible can be proven by a contradiction. Suppose A^r , B^r , and C^r are compatible, so that there exists a joint POVM G with elements G_{ijk} such that $A_i^r = \sum_{jk} G_{ijk}$, $B_j^r = \sum_{ik} G_{ijk}$, and $C_k^r = \sum_{ij} G_{ijk}$. As $\langle \varphi_i | A_i^r \varphi_i \rangle = 0$ and all the operators G_{ijk} are positive, we have $\langle \varphi_i | G_{ijk} \varphi_i \rangle = 0$ and this further implies $G_{ijk} \varphi_i = 0$. Similarly, $G_{ijk} \psi_j = 0$ and $G_{ijk} \chi_k = 0$. Since the set $\{\varphi_i, \psi_j, \chi_k\}$ spans \mathcal{H}_3 , we conclude that $G_{ijk} = 0$. This contradicts the normalization $\sum_{ijk} G_{ijk} = 1$. Hence, the three POVMs A^r , B^r , and C^r are incompatible.

The sufficient condition (18) for compatibility of the reverse versions of regular rank-1 POVMs is not necessary. We next demonstrate that there are compatible observables that do not satisfy (18).

Example 8: Two mutually unbiased bases. Consider a d-dimensional Hilbert space \mathcal{H}_d and an orthonormal basis $\{\varphi_i\}_{i=0}^{d-1}$ in it. We denote $\omega=e^{i2\pi/d}$ and define another

orthonormal basis $\{\psi_j\}_{j=0}^{d-1}$ by

$$\psi_j = \frac{1}{\sqrt{d}} \sum_{k=0}^{d-1} \omega^{jk} \varphi_k. \tag{19}$$

These two bases are mutually unbiased, meaning that $|\langle \varphi_i | \psi_j \rangle| = \frac{1}{\sqrt{d}}$ for all $i, j = 0, \dots, d-1$. The related POVMs $A_i = |\varphi_i\rangle\langle\varphi_i|$ and $B_j = |\psi_j\rangle\langle\psi_j|$ consist of noncommuting projections and are hence incompatible.

The reverse versions A^r and B^r are incompatible if d=2, since then A^r and B^r are just relabelings of A and B. However, for any $d \ge 3$, A^r and B^r are compatible even if inequality (18) does not hold. To see this, we recall that, by Proposition 2 in [39], A^r and B^r are compatible whenever there exists a quantum state $\sigma \in \mathcal{S}(\mathcal{H})$ such that

$$tr[\mathsf{A}_i\sigma] = \frac{1 - \delta_{i0}}{d - 1} \quad \text{and} \quad tr[\mathsf{B}_j\sigma] = \frac{1 - \delta_{j0}}{d - 1}. \tag{20}$$

It is not hard to check that the operator

$$\sigma = \frac{1}{d-1} \sum_{i=1}^{d-1} |\varphi_i\rangle \langle \varphi_i|$$
$$-\frac{1}{(d-1)(d-2)} \sum_{1 \le i < j \le d} (|\varphi_i\rangle \langle \varphi_j| + |\varphi_j\rangle \langle \varphi_i|)$$

is a density operator and satisfies the conditions above. Therefore, A^r and B^r are compatible.

As explained in Example 6, the reversing channel v^r can also be applied to an already reverse observable A^r to obtain a doubly reverse observable A^{rr} . It is not hard to see from Proposition 1 that two doubly reverse observables are always compatible if their number of outcomes $N \ge 3$. More generally, a sufficient condition for compatibility of m doubly reverse observables with N outcomes each is $m \le (N-1)^2$.

B. Eigenvalue condition for PPOVMs

Let A be a PPOVM with an outcome set X and the normalization $\sum_{x \in X} A_x = \varrho \otimes \mathbb{1}$ for some state ϱ . We denote by m_x the minimal eigenvalue of the PPOVM element A_x for each $x \in X$. The noise content of A satisfies

$$w(\mathsf{A};\mathcal{T})\geqslant\sum_{x\in\mathcal{X}}m_{x}.$$
 (21)

To see this, we define a trivial PPOVM T as

$$\mathsf{T}_{x} = \frac{m_{x}}{m} \varrho \otimes \mathbb{1},\tag{22}$$

where $m = \sum_{x \in X} m_x$. Since

$$\mathsf{A}_x \geqslant m_x \mathbb{1} \otimes \mathbb{1} \geqslant m_x \varrho \otimes \mathbb{1},\tag{23}$$

we can define

$$\mathsf{A}_x' = \frac{1}{1-m} (\mathsf{A}_x - m_x \varrho \otimes \mathbb{1}) \tag{24}$$

and A' is a valid PPOVM. We can then write

$$A = mT + (1 - m)A', \tag{25}$$

which confirms (21). Proposition 4 thus implies the following result, analogous to Corollary 1.

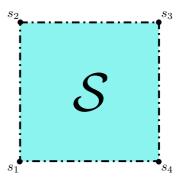


FIG. 5. Squit state space.

Corollary 2. If $A^{(1)}, \dots, A^{(m)}$ is a collection of m incompatible PPOVMs, then the sum of the minimal eigenvalues of all their effects is smaller than m-1.

We note that, in contrast to the case of POVMs, the eigenvalue formula (21) provides only a lower bound for the noise content of a PPOVM. For instance, let

$$\mathsf{A}_{x} = p_{x} |\psi_{x}\rangle \langle \psi_{x}| \otimes \mathbb{1},\tag{26}$$

where $\langle \psi_x | \psi_y \rangle = \delta_{xy}$ and p_x is a probability distribution. Then $m_x = 0$ for all x and the right-hand side of (21) equals zero. But the PPOVM A is trivial, so the left-hand side of (21) equals 1.

C. Polytope state spaces

A compact convex subset P of a finite-dimensional vector space V is a *polytope* if it has a finite number of extreme elements. Let $ext(P) = \{s_1, \ldots, s_n\}$ be the set of extreme elements of a polytope P. Since every state $s \in P$ can be represented as a convex sum of elements in ext(P), we have that

$$\mathsf{A}_{x}(s) = \mathsf{A}_{x}\bigg(\sum_{i} \lambda_{i} s_{i}\bigg) = \sum_{i} \lambda_{i} \mathsf{A}_{x}(s_{i})$$

$$\geqslant \sum_{i} \lambda_{i} \min_{k} \mathsf{A}_{x}(s_{k}) = \min_{k} \mathsf{A}_{x}(s_{k})$$

for every $s \in P$, and thus $\inf_{s \in S} A_x(s) = \min_{s \in ext(P)} A_x(s)$. Combining this with Proposition 4, we get an analogous result to the previous eigenvalue conditions for POVMs and PPOVMs.

Corollary 3. If $A^{(1)}, \ldots, A^{(m)}$ is a collection of m incompatible observables on a polytopic state space P, then the sum of minimal values of all of their effects on ext(P) is smaller than m-1

In the following, we take S to be a state space that is isomorphic to a square in \mathbb{R}^2 , i.e., to the convex hull of four points $s_1, s_2, s_3, s_4 \in \mathbb{R}^2$ satisfying $s_1 + s_3 = s_2 + s_4$ (see Fig. 5). This is called the square bit state space, or squit state space for short.

We consider a class of binary observables A^{α} and B^{β} , parametrized by $\alpha, \beta \in [0,1]$, whose outcomes are labeled by \pm and defined on the extreme points s_1 , s_2 , s_3 , and s_4 as

$$A_+^{\alpha}(s_1) = A_+^{\alpha}(s_2) = \alpha, \quad A_+^{\alpha}(s_3) = A_+^{\alpha}(s_4) = 1,$$

$$\mathsf{B}_{+}^{\beta}(s_1) = \mathsf{B}_{+}^{\beta}(s_4) = \beta, \quad \mathsf{B}_{+}^{\beta}(s_2) = \mathsf{B}_{+}^{\beta}(s_3) = 1.$$

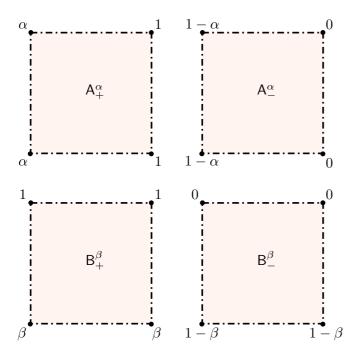


FIG. 6. Observables A^{α} and B^{β} .

The values of A^{α} and B^{β} are depicted in Fig. 6. Now we see that

$$w(\mathsf{A}^{\alpha}; \mathcal{T}) = \min_{s \in \mathsf{ext}(S)} \mathsf{A}_{+}(s) + \min_{s \in \mathsf{ext}(S)} \mathsf{A}_{-}(s) = \alpha,$$

and similarly that $w(\mathsf{B}^\beta;\mathcal{T})=\beta$. Hence, by Corollary 3, if

$$\alpha + \beta \geqslant 1,\tag{27}$$

then observables A^{α} and B^{β} are compatible. It is easy to find A^{α} and B^{β} as mixtures with maximal noise contents,

$$A^{\alpha} = \alpha T + (1 - \alpha)A,$$

$$B^{\beta} = \beta T + (1 - \beta)B,$$

where T is the trivial binary observable with $T_+(s) = 1$ and $T_-(s) = 0$ for all $s \in S$, and $A \equiv A^0$ and $B \equiv B^0$.

The observables A and B are themselves incompatible. More so, they are maximally incompatible in the sense that the minimum amount of noise one has to mix them with to make their noisy versions compatible is enough to make any other pair of observables compatible. More precisely, it was shown in [21] that the observables $\lambda A + (1-\lambda)T_1$ and $\mu B + (1-\mu)T_2$ are incompatible for all choices of trivial observables T_1 and T_2 if and only if $\lambda + \mu > 1$. Therefore, we conclude that inequality (27) derived from Proposition 2 is actually both necessary and sufficient for the compatibility of A^{α} and B^{β} .

V. CONCLUSIONS

We have considered general probabilistic theories on an equal footing and quantified the noise content of observables in every such theory via the set of trivial observables. In the case of standard quantum theory, the noise content is merely the sum of minimal eigenvalues of the POVM effects. In the quantum theory of processes, the noise content is bounded below by the sum of minimal eigenvalues of the corresponding PPOVM effects. In general, the noise content can be quantified with respect to any subset of observables.

We have derived the noise content inequality for a pair of observables, which is a necessary condition for their incompatibility. Our approach is based on a modification of the adaptive strategy for building a joint observable. We have then extended this result to the case of *m* observables. By way of examples with reverse regular observables we have demonstrated nontriviality of the derived noise content inequality. Moreover, this inequality turned out to be not only necessary but also sufficient for incompatibility of some observables in the square bit state space.

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